

**MATH455 HOMEWORK 7**  
**DUE FRIDAY, APRIL 3**

Do Exercise 1 and your choice of one of Exercise 2, 3, or 4.

*Exercise 1.* Let  $H'$  be the set of all pairs  $(P, \vec{R})$  of register machine programs  $P$  and input  $\vec{R}$  so that  $P$  halts with  $\vec{R}$  as input. Argue directly that  $H'$  is not decidable. (That is, do not derive this as a corollary of the version of the halting problem from lecture. Instead, use a similar diagonal argument to show that  $H'$  is not decidable.)

*Exercise 2.* Consider the following “proof” that there is no computable listing of register machine programs:

- Suppose toward a contradiction that there is a computable listing of register machine programs so that each register machine is assigned an index  $e$  so that from  $e$  we can compute the corresponding program and vice versa. Now define a register machine program  $B$ , taking input a natural number  $e$ . Running  $B$  on input  $e$  does the following: First as a subroutine  $B$  runs the  $e$ -th register machine program on input  $e$ , getting some output  $y \in \mathbb{N}$ . This can be done similar to how we defined a universal register machine. Then  $B$  takes this subroutine output  $y$  outputs  $y + 1$ .

Fix the index  $e$  for the program  $B$ . Consider what happens when  $B$  is run with  $e$  as input. It simulates running  $B$  with input  $e$ , producing some output  $y \in \mathbb{N}$ . But by the definition of  $B$  we get that it outputs  $y + 1$ . Since both  $y$  and  $y + 1$  are the output to  $B$  with the same input  $e$  we get that  $y = y + 1$ , which is impossible. This contradiction shows that our original assumption that there is a computable listing of register machine programs must be false.

Identify the flaw in this “proof” and explain why it is not valid.

*Exercise 3.* Prove that for any natural number  $e$  there are infinitely many natural numbers  $e'$  so that  $\varphi_e = \varphi_{e'}$ .<sup>1</sup>

*Exercise 4.* Prove that for each natural number  $s$  the set  $\{P : P \text{ is a register machine program which halts in } \leq s \text{ many steps}\}$  is decidable.

---

<sup>1</sup>Recall from the video for Friday’s lecture that  $\varphi_e$  is the function computed by the  $e$ -th register machine program, according to some fixed computable listing of the register machine programs.