#### Nonstandard methods versus Nash-Williams

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# Our project

- $\bullet$  Nonstandard methods have been fruitfully applied to prove theorems about combinatorics on  $\mathbb N$ 
  - Namedrop: Di Nasso, Goldbring, Jin, Tao, ...
- Topological Ramsey theory studies combinatorial topological spaces which generalize Ellentuck space ( $\approx$  the space of subsets of  $\mathbb{N}$ ), the familiar setting for ordinary Ramsey theory
- Let's apply nonstandard methods to a more general setting than Ellentuck space
- Starting point: the Nash-Williams theorem for Ellentuck space and its generalization

We can use tools from model theory to prove theorems outside of logic

- Take a structure. For this talk, it will mostly be N
- Take an ultrapower of N to embed N into a saturated elementary extension \*N
- Exploit the connection  $\mathbb{N} \hookrightarrow {}^*\mathbb{N}$  to prove theorems about  $\mathbb{N}$



# A gentle warmup: the pigeonhole principle

#### Theorem (Pigeonhole Principle)

If you partition  $\mathbb{N}$  into finitely many pieces  $X_0, \ldots, X_n$  then one of the pieces is infinite.

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#### **Proof:**

- Consider  $\alpha \in {}^*\mathbb{N} \setminus \mathbb{N}$
- \*X<sub>0</sub>,...,\*X<sub>n</sub> are a partition of \*N (by elementarity)
- So  $\alpha$  is in some  $*X_i$
- So X<sub>i</sub> is infinite (by elementarity)



# Iterating the \* map

I lied earlier when I said nonstandard methods work by embedding  $\mathbb N$  into  ${}^*\mathbb N$ 

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- Actually we embed  $V_{\omega}(\mathbb{N})$  into a saturated elementary extension
- Then  $^*V_{\omega}(\mathbb{N})$  is a definable class in  $V_{\omega}(\mathbb{N})$
- So <sup>\*</sup>ℕ is a set in the domain of the embedding
- We can apply the \* map to it and its elements
- If  $\alpha \in {}^*\mathbb{N} \setminus \mathbb{N}$  then  $\alpha < {}^*\alpha$
- And we can iterate:

$$\mathbb{N} \hookrightarrow {}^*\mathbb{N} \hookrightarrow {}^{*(2)}\mathbb{N} \hookrightarrow \cdots \hookrightarrow {}^{*(k)}\mathbb{N} \hookrightarrow \cdots$$



#### Theorem (Ramsey 1930)

Partition  $[\mathbb{N}]^k$  into finitely many pieces  $X_0, \ldots, X_n$ . Then there is infinite  $H \subseteq \mathbb{N}$  so that  $[H]^k \subseteq X_i$  for some *i*.

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Proof (k = 3):

- Consider  $\alpha \in *\mathbb{N} \setminus \mathbb{N}$
- Then  $\langle \alpha, {}^{*}\alpha, {}^{*(2)}\alpha \rangle$  is in some  ${}^{*(3)}X_i$
- So  $\alpha \in {}^{*}{\{a \in \mathbb{N} : \langle a, \alpha, {}^{*}\alpha \rangle \in {}^{*(2)}X_i\}}.$
- So {a ∈ N : ⟨a, α, \*α⟩ ∈ \*(2)X<sub>i</sub>} is infinite
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Now induct:

- Already built  $H_i = \langle h_0, \ldots, h_i \rangle$
- Inductively,  $lpha \in {}^{*}\{a \in \mathbb{N} : t^{\frown}a \in X_i\}$  for each  $t \in [H_i]^2$
- And  $\alpha \in {}^{*}{a \in \mathbb{N} : t^{\alpha}a^{\alpha} \in {}^{*}X_{i}}$  for each  $t \in [H_{i}]^{1}$

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- $\bullet$  Finitely many, and  $\alpha$  is in their nonstandard intersection
- So their standard intersection is infinite
- Pick  $h_{i+1} > h_i$  from that intersection

Finally  $H = \langle h_i \rangle$  is monochromatic

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# Generalizing Ramsey to families of sets of nonuniform size

#### Definition

The Schreier barrier S consists of all  $s \in [\mathbb{N}]^{<\omega}$  so that  $|s| = \min s + 1$ .

- The first element of *s* tells you how long *s* is
- You can think of S as a tagged amalgamation of (copies of) all [ℕ]<sup>k</sup>



# A Ramsey property for the Schreier barrier

#### Theorem (Nash-Williams for S)

Partition S into finitely many pieces  $X_0, \ldots, X_n$ . Then there is infinite  $H \subseteq \mathbb{N}$  so that  $S \upharpoonright H$  is monochromatic.

$$\mathcal{S} \upharpoonright \mathcal{H} = \{ s \in \mathcal{S} : s \subseteq \mathcal{H} \}$$
  
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- For  $[\mathbb{N}]^k$  we looked at what piece of the partition contained  $\langle \alpha, {}^*\alpha, \dots, {}^{*(k-1)}\alpha \rangle$
- But now we don't know in advance how long a sequence in S will be
- Intuitively, we want to look at

$$\langle \alpha, {}^*\alpha, \dots {}^{*(\alpha)}\alpha \rangle$$

• But this is nonsensical—what would it even mean to iterate \* a nonstandard number of times?

# A proxy for $\langle \alpha, {}^*\alpha, \dots {}^{*(\alpha)}\alpha \rangle$

Notation:

• \*
$$\mathbb{N} = \dim_{k \in \omega} {}^{*(k)}\mathbb{N}$$

• For  $\beta \in *\mathbb{N}$ , let  $k(\beta)$  be the least k so that  $\beta \in {}^{*(k)}\mathbb{N}$ 

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• For  $\beta \in *\mathbb{N}$ , let  $k(\beta)$  be the least k so that  $\beta \in {}^{*(k)}\mathbb{N}$ 

**Claim:** Fix  $\alpha \in *\mathbb{N}$ . For any sequence  $\langle \beta_i : i \in \omega \rangle$  there is (a non-unique)  $\sum_{\alpha} \beta_i \in *\mathbb{N}$  so that for all  $X \subseteq \mathbb{N}$ 

$$\sum_{i\in\mathbb{N};\,\alpha}\beta_i\in{}^*X\quad\Leftrightarrow\quad\alpha\in{}^*\{i\in\mathbb{N}:\beta_i\in{}^{*(k(\beta_i))}X\}$$

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 $\bullet$  Our proxy for  $\langle \alpha, {}^{*}\alpha, \ldots, {}^{*(\alpha)}\alpha\rangle$  is then

$$\sigma(\alpha) = \sum_{i \in \mathbb{N}; \alpha} \langle \alpha, \dots, {}^{*(i)} \alpha \rangle$$

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$$S \upharpoonright H = \{ s \in S : s \subseteq H \}$$
  

$$s_k = \langle \alpha, \dots *^{(k)} \alpha \rangle \text{ approximate } \sigma(\alpha)$$
  
**Proof:**

- Consider  $\alpha \in {}^*\mathbb{N} \setminus \mathbb{N}$
- Then  $\sigma(\alpha)$  is in some  $*X_i$
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# Further generalization: fronts

- $\mathcal{F} \subseteq [\mathbb{N}]^{<\omega}$  is a front if
  - (antichain or Nash-Williams property)  $s \not\sqsubseteq t$  for  $s \neq t$  from  $\mathcal{F}$
  - (density)

For any infinite  $b \subseteq \mathbb{N}$  there is  $s \sqsubseteq b$  from  $\mathcal{F}$ 

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Examples:

- $[\mathbb{N}]^k$  for any k
- The Schreier barrier  ${\cal S}$

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To prove a Ramsey property for  $[\mathbb{N}]^k$  and S we had an idea of what a generic nonstandard member looked like, based on how the front was built up

- $\langle \alpha, \dots, {}^{*(k-1)}\alpha \rangle$  for  $[\mathbb{N}]^k$
- $\sigma(\alpha)$ , a proxy for  $\langle lpha, \dots, {}^{*(lpha)} lpha 
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If we want to do the same for an arbitrary front  ${\cal F}$  we need to understand how  ${\cal F}$  is built up

### Trees of fronts

For  ${\mathcal F}$  a front, set

 $\mathcal{T}(\mathcal{F}) = \{t \in [\mathbb{N}]^{<\omega} : t \sqsubseteq s \text{ for some } s \in \mathcal{F}\}$ 

Then  $T(\mathcal{F})$  is a tree and  $\mathcal{F}$  are the leaves

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- If b were an infinite branch through T(F) it'd extend some s ∈ F by density
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We can think of  ${\mathcal F}$  as built up by induction on  ${\mathcal T}({\mathcal F})$ 

- For  $s \in \mathcal{F}$ , set  $\mathcal{F}_s = \{s\}$
- ullet For  $s\in \mathcal{T}(\mathcal{F})\setminus \mathcal{F}$ , set  $\mathcal{F}_s=igcup_{t\in ext{succ}\,s}\mathcal{F}_t$
- Here succ s is the set of successors of s in  $T(\mathcal{F})$
- Observe that  $\mathcal{F}_s$  is a front on  $[\mathbb{N}]^{<\omega} \upharpoonright s$

Finally  $\mathcal{F}=\mathcal{F}_{\emptyset}$ 

$$\mathcal{S} = \{ s \in [\mathbb{N}]^{<\omega} : |s| = \min s + 1 \}$$

What is  $S_s$  for subsequences s of (2,7,9)?

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•  $S_{\langle 2,7 \rangle} = \{ \langle 2,b,c \rangle : 2 < b < c \} = \{ 2^{\uparrow} t \}$ 

$$\begin{array}{l} \mathcal{C}_{2} = \{(2, b, c) : 2 < b < c\} = \{2, c, c\} \\ t \in [\mathbb{N} \setminus 3]^2 \} \end{array}$$

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$$\mathcal{S}_{\langle 2 \rangle} = \{ \langle 2, b, c \rangle : 2 < b < c \} = \{2^{\uparrow}t : t \in [\mathbb{N} \setminus 3]^2 \}$$

• 
$$\mathcal{S} = \mathcal{S}_{\emptyset} = \{a^{\frown}t : t \in [\mathbb{N} \setminus (a+1)]^a\}$$

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#### Theorem (Nash-Williams theorem)

Let  $\mathcal{F}$  be a front. Partition  $\mathcal{F}$  into finitely many pieces  $X_0, \ldots, X_n$ . Then there is infinite  $H \subseteq \mathbb{N}$  so that  $\mathcal{F} \upharpoonright H$  is monochromatic.

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, set  $\sigma_s = \sigma_s(lpha)$  to be  $\langle lpha 
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, set  $\sigma_s = \sigma_s(\alpha)$  to be  $\sum_{t \in \text{succ } s; \alpha} \sigma_t(\alpha)$ 

Recall:

$$\sum \sigma_t \in {}^*X \quad \Leftrightarrow \quad s^{\frown} \alpha \in {}^*\{a \in \mathbb{N} : \sigma_{s^\frown a} \in {}^*X\}$$

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### The Nash-Williams theorem for Ellentuck space

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- $\sigma_{\emptyset}(\alpha)$  is in some  $^{*}X_{i}$
- Pick  $h_0$  to be the minimum element of  $\{a \in \mathbb{N} : \sigma_a \in {}^*X_i\}$
- Then inductively pick  $h_{i+1} > h_i$  using that  $\alpha$  is in  $\{a \in \mathbb{N} : \sigma_{t \cap a} \in X_i\}$  for each subset t of the *i*-th partial solution  $H_i$

Finally  $H = \langle h_i \rangle$  is monochromatic

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Ellentuck space  $\mathcal{E}$  has multiple components

- The points are elements of  $[\mathbb{N}]^\omega$
- You can associate to each point its k-th finite approximation in [ℕ]<sup>k</sup>
- $\bullet$  There is a partial order  $\subseteq$  on points

### Abstract Ramsey spaces

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And  $\mathcal{E}$  has some nice properties

- (A.1) Sequencing: points behave like infinite sequences
- (A.2) Finitization: you can port the partial order ⊆ to the finite approximations, and each approximation has a finite number of predecessors
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A Ramsey space is a tuple  $(\mathcal{R}, \mathcal{AR}, \leq, r)$  satisfying (A.1–4) where  $\mathcal{R}$  are the points,  $r : \mathcal{R} \times \omega \to \mathcal{AR}$  is the finite approximation map, and  $\leq$  is the partial order

# The topological in topological Ramsey theory

The Ellentuck topology on  ${\mathcal R}$  is generated by basic open sets

 $[a,X] = \{Y \in \mathcal{R} : Y \leq X \text{ and } \exists k r_k(Y) = a\}.$ 

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If  $\mathcal{R}$  is closed as a subspace of the product topology on  $\mathcal{AR}$ , it's quite nice

- X ⊆ R is Ramsey if you can refine any basic open set be either contained in or disjoint from X
- *X* ⊆ *R* is Ramsey null if it is Ramsey and you can always refine to be disjoint from *X*

- If  $\mathcal{R}$  is closed, any Baire subset is Ramsey and any meager subset is Ramsey null
- Indeed any Souslin-measurable or Borel subset is Ramsey

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#### The abstract Nash-Williams theorem

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#### Theorem (Abstract Nash-Williams)

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- I'd like to say our nonstandard proof of the Nash-Williams theorem extends directly to the full abstract Nash-Williams theorem
- But we need the space to be amenable to nonstandard methods
- And we don't (yet?) have a proof that this applies to every nontrivial Ramsey space

### What we do have for the abstract Nash-Williams theorem

Under an extra assumption the nonstandard proof goes through.

Theorem (Partial abstract Nash-Williams)

Consider a front  $\mathcal{F}$  on  $\mathcal{AR}$ . Suppose

- $\mathcal{AR}$  is infinitely branching everywhere; and
- There is a filter C on  $\mathcal{R}$  so that for each  $s \in T(\mathcal{F}) \setminus \mathcal{F}$  the restriction of succ s to C is a nonprincipal ultrafilter on succ s.

Then  $\mathcal{F}$  satisfies a Ramsey partition property.

- $\bullet~(\mathcal{R},\leq)$  is a poset, so the usual definition of filter applies to  $\mathcal C$
- succ  $s \upharpoonright X = \{t \in \operatorname{succ} s : \exists k \ t \leq_{\operatorname{fin}} r_k(X)\}$
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Any Ramsey space which can be thought of as its (k + 1)-th approximations coming from k-th approximations by concatenating sequences from (cofinite subsets of) a countable alphabet will admit such a filter

- Ellentuck space
  - Restrict any nonprincipal ultrafilter on  $\mathcal{P}(\mathbb{N})$  to the infinite subsets to get  $\mathcal C$

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- Ellentuck space
  - Restrict any nonprincipal ultrafilter on  $\mathcal{P}(\mathbb{N})$  to the infinite subsets to get  $\mathcal C$
- The Milliken space of block sequences
- The Hales-Jewett space of variable words
- The space  $\mathcal{E}_{\omega}(\mathbb{N})$  of equivalence relations on  $\mathbb{N}$  with infinite quotients

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# A silly negative example

#### The V space

- V has two points, the constant 0 sequence and the constant 1 sequence
- Finite approximations are finite constant 0 or 1 sequences
- $\bullet$  Trivially, any front on  $\mathcal{AV}$  satisfies a Ramsey partition property
- But  ${\mathcal V}$  doesn't satisfy the filter property!



# Continuing work

#### Question

Suppose you have a nontrivial<sup>a</sup> Ramsey space  $(\mathcal{R}, \mathcal{AR}, \leq, r)$  and a front  $\mathcal{F}$  on  $\mathcal{AR}$ . Then there is a filter  $\mathcal{C}$  on  $\mathcal{R}$  so that for each  $s \in T(\mathcal{F}) \setminus \mathcal{F}$  the restriction of succ s to  $\mathcal{C}$  is a nonprincipal ultrafilter on succ s.

<sup>a</sup>What should this mean?

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- The abstract Nash-Williams theorem isn't the only theorem in abstract Ramsey theory
- What other results are amenable to nonstandard methods?

# Thank you!

K. Williams (SHSU  $\rightarrow$  SR)

Nonstandard Methods vs. Nash-Williams

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